Computationally, Tone is Different

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Introduction

- Hyman (2011) asks: Is tone different?
- Automata-theoretic measure of complexity also says yes
- Typological studies provide evidence segmental phonology is *regular* and *subsequential* (Gainor et al. (2012), Chandlee and Heinz (2012), Heinz and Lai (2013), Chandlee, (in prep.))
- A common tonal process, *Unbounded Tonal Plateauing* (Hyman, 2011, henceforth UTP), is *regular*, but not *subsequential*, and thus more complex
- Gives us intuition that *unbounded*, *bidirectional* processes are not expected for segmental phonology

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Why tone?

- Hyman (2011) argues that tone "can do everything that segmental or metrical phonology can do, but the reverse is not true" (p.236)
- Yip (2002) lists characteristics uniquely common in tone: ex. 'mobility' and 'one-to-many'
- An example from Digo (Bantu; Kisseberth, 1984)
 - a. ni-na+tsukur-a 'l am taking'
 b. ni-na+<u>a</u>-tsukŭr-â 'l am taking them'

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Why tone?

- Is tone more *complex* than segmental phonology?
- Segmental phonology also has long-distance and 'one-to-many' processes
- Can computational measures of complexity give us a hard distinction?

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• Much of phonology deals with change from underlying representation (UR) to surface representation (SR)

(2)
$$n \rightarrow m / _ p$$

• The change (2) can be modeled as a string-to-string mapping

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An example

• *Regular* string-to-string mappings can be modeled with finite-state *transducers* (FSTs)

(4)	A FST for n $ ightarrow$ m / p	



CV	np⊢→C	Vmp
Input	0	Output
С	0	C
V	0	V
n	1	
р	0	mp

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Why automata?

(4) n
$$\rightarrow$$
 m /___ p



- Can look at binary UR/SR relation independent of phonological theory
- (4) is a *subsequential* FST (SFST)
- At each state, one transition per input

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Complexity of mappings

- SFSTs are *strictly less expressive* than FSTs Mohri (1997)
- Mappings describable with SFSTs fall into two subsequential subregions of the *regular* region
- Typological studies of segmental processes show they fall into one of these regions



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Complexity of mappings

• Even long-distance regressive harmony is subsequential

 $(5)...[-F]...[+F]... \ \mapsto ...[+F]...[+F]...$



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Complexity of mappings

- Subsequentiality is shared by local and long-distance segmental processes
- UTP, a tonal process, is not

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Unbounded Tone Plateauing

- UTP (Hyman, 2011) is a tonal process in which all TBUs inbetween two H tones become H (assuming underlying H/\emptyset)
 - (6) Luganda (Bantu; Hyman et al., 1987; Hyman and Katamba, 2010)
 - a. /bikopo/ \rightarrow bikópo $\varnothing H \varnothing \rightarrow L H L$ 'cups'
 - b. /byaa-walusiimbi/ \rightarrow byaa-walúsiimbi $\varnothing \ \varnothing \ H \varnothing \ \varnothing \ \rightarrow \ L \ L \ H \ L \ L$ 'of Walusimbi'
 - c. /bikopo byaa-walusiimbi/ \rightarrow bikópó byáá-wálúsiimbi $\varnothing H \varnothing \ \varnothing \ \Theta H \varnothing \ \varnothing \rightarrow \ L H H H H H L L$ 'the cups of Walusimbi'

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Unbounded Tone Plateauing

• Can be formalized as the following:

(7) $\varnothing^n \to \mathsf{H}^n / \mathsf{H} _ \mathsf{H}$

• Attested in Luganda (Hyman and Katamba, 2010; Hyman et al., 1987), Zulu (Cassimjee and Kisseberth, 2001), Kihunde (Goldsmith, 1990), Amahuaca (Russel and Russel, 1959), and others

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UTP as a mapping



 $\varnothing \varnothing \varnothing \mapsto \mathsf{LLL}$

 $\varnothing H \varnothing \mapsto \mathsf{LHL}$

 $\varnothing H \varnothing H \varnothing \mapsto L H H H L$

- This machine is *not* a SFST
- I have a formal proof, based on properties of subsequentiality given in Oncina et al. (1993), showing we *cannot* build an SFST for (7)

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UTP as a non-subsequential mapping

 Recall that SFSTs can 'wait' some set time before writing an output



 $\varnothing \varnothing \mapsto \mathsf{LLL}$ $\varnothing \mathsf{H} \varnothing \mathsf{H} \varnothing \mapsto \mathsf{LHHHL}$ $\mathsf{H} \varnothing \varnothing \mathsf{H} \mapsto \mathsf{*HLLLH}$

(8) SFST with one wait state

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UTP as a non-subsequential mapping

 Recall that SFSTs can 'wait' some set time before writing an output



 $\varnothing \varnothing \mapsto LLL$ $\varnothing H \varnothing H \varnothing \mapsto LHHHL$ $H \varnothing \varnothing H \mapsto HHHH$ $H \varnothing \varnothing \vartheta H \mapsto *HLLLH$

(9) SFST with two wait states

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UTP as a non-subsequential mapping

 Recall that SFSTs can 'wait' some set time before writing an output



(10) SFST with three wait states

 $\begin{array}{l} \mathsf{H} \varnothing \varnothing \varphi \mathsf{H} \mapsto \mathsf{H} \mathsf{H} \mathsf{H} \mathsf{H} \mathsf{H} \\ \mathsf{H} \varnothing \varnothing \varnothing \mathsf{H} \mapsto \mathsf{*} \mathsf{H} \mathsf{L} \mathsf{L} \mathsf{L} \mathsf{L} \mathsf{H} \end{array}$

- A SFST with n number of wait states will fail when Hs are n + 1 TBUs apart
 - To capture UTP, an SFST needs an infinite number of wait states

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The complexity of UTP

- UTP cannot be modeled with a SFST
- It is not a subsequential mapping
- This comes from the *unbounded*, *bidirectional* nature of the mapping
- Unlike harmony, there are *two* triggers on either side of the target, arbitrarily far away
- It *can* be modeled with a non-subsequential FST, so it is regular

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Conclusions



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Conclusions

- UTP, a common tonal process, is not subsequential
- Tone *is* different; it appears to be more computationally complex
- New generalization: tone can have *unbounded*, *bidirectional* processes; segmental processes cannot
- There are two potential exceptions to this: Sanskrit *nati* (Whitney, 1889; Macdonell, 1910), and Yaka vowel 'plateauing' harmony (Hyman, 1998), a VH version of UTP
- Subsequentiality of segmental processes remains robust, but we can be on the look out for unbounded, bidirectional processes

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