Computational Locality and Autosegmental Processes

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Introduction

- Autosegmental representations (ARs) (Goldsmith, 1976; Clements, 1976) have been claimed to capture non-local phenomena in a local way (McCarthy, 1982; Odden, 1994).
- We apply a computational notion of locality to a selection of tone processes to get a more nuanced understanding of this ability of ARs.
- Three-way distinction:
 - Local even without ARs.
 - Local only with ARs.
 - Not local even with ARs.

Computational notion of locality

- Based on the Input Strictly Local (ISL) functions, which were originally defined in terms of formal language theory and automata theory (Chandlee, 2014).
- We'll be using the logical characterization of ISL proposed by Chandlee and Lindell (to appear).
 - ▶ ISL function = Quantifier-free First Order Graph Interpretation
- Why use logic?
 - We can directly extend a restrictive, explicit notion of locality from strings to phonological representations

 (1) Rimi (Schadeberg, 1979; Meyers, 1997) /u-pµm-a/ → [u-pµm-á] 'to go away'



UR string model defined with:

- 5 positions, labeled with segments
- ▶ successor function: s(0) = 1, s(1) = 2, ..., s(4)=4
- ▶ predecessor function p(4) = 3, p(3) = 2, ..., p(0)=0

The SR string graph is defined in terms of the UR graph using FO logic formulas (Engelfriet and Hoogeboom, 2001).



- ► *H*(*x*) is True iff position *x* bears a high tone.
- V(x) is True iff position x is a vowel.
- $\varphi_{\acute{V}} \stackrel{\mathsf{def}}{=} V(x) \wedge H(p(p(x)))$
 - An output position bears a high tone iff it's a vowel and the previous vowel bears a high tone in the input graph.

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ISL = Quantifier-free FO Logic

$$\varphi_{\acute{V}} \stackrel{\mathsf{def}}{=} V(x) \wedge H(p(p(x)))$$

- These formulas do not use the full power of FO: they don't use quantifiers.
- The processes that can be described in this way are those for which the trigger and the target form a contiguous substring of bounded length in the input string.
- The boundedness means we can use the successor or predecessor function repeatedly to determine whether both the target and trigger are present.
 - No quantifier is needed. (e.g, no "... ∧ (∃z)[...]")

Autosegmental Representation (AR) Graphs

 Goal: extend this same notion of locality (QF FO describable) from string graphs to AR graphs (Jardine, 2016).

Example: VÚVÚÚV



Autosegmental Representation (AR) Graphs



 $\begin{smallmatrix} H & H \\ | & | \\ V & V & V & V \end{smallmatrix}$

Method

- We designate as ISL those patterns that can be describe with QF FO using string graphs.
- We designate as AISL those patterns that can be described with QF FO using AR graphs.
- We illustrate that ISL patterns are also AISL but not vice versa
- ▶ We will also identify cases that are neither ISL nor AISL.

Process		ISL?	AISL?
Bounded spread	(Bemba)	\checkmark	
Bounded shift	(Rimi)	\checkmark	
Unbounded shift	(Zigula)	X	
Unbounded OCP	(Arusa)	X	
Unbounded spread	(Ndebele)	X	
Meussen's rule	(Shona)	X	

Bounded spread

 ▶ Bemba (Bickmore and Kula, 2013) /bá-la-kak-a/ → [bá-lá-kak-a] 'they tie up' /tu-la-bá-kak-a/ → [tu-la-bá-kák-a] 'we tie them up'

$$V \lor V \lor V \lor V \lor V \lor V \lor V$$

Bounded spread

$$a_O(x, \mathbf{y}) \stackrel{\text{def}}{=} a_I(x, \mathbf{y}) \lor a_I(p(x), \mathbf{y})$$

$$\bigvee_{V} \bigvee_{V} \bigvee_{V$$

Bounded shift

► Rimi (Schadeberg, 1979; Meyers, 1997) /u-pým-a/ \mapsto [u-pým-á] 'to go away' /rá-mu-ntu/ \mapsto [ra-mú-ntu] 'of a person' /mu-tém-j/ \mapsto [mu-tem-j] 'chief' H H H V V V V V V V V V V V V

Bounded shift

$$a_O(x, y) \stackrel{\mathsf{def}}{=} a_I(p(x), y)$$

$$V \lor V \lor V \lor V \lor V \lor V \lor V$$

Process		ISL?	AISL?
Bounded spread	(Bemba)	\checkmark	\checkmark
Bounded shift	(Rimi)	\checkmark	\checkmark
Unbounded shift	(Zigula)	X	
Unbounded OCP	(Arusa)	X	
Unbounded spread	(Ndebele)	X	
Meussen's rule	(Shona)	X	

Unbounded shift

 Zigula (Kenstowicz and Kisseberth, 1990) ku-gulus-a 'to chase' ku-lombéz-a 'to ask' ku-lombez-éz-a 'to ask for' ku-lombez-ez-án-a 'to ask for each other'

Unbounded shift

$$\begin{aligned} \texttt{lastH}(y) &\stackrel{\text{def}}{=} H(y) \land s(y) = y \\ \texttt{penultV}(x) &\stackrel{\text{def}}{=} V(x) \land (s(s(x)) = s(s(s(x))) \\ a_O(x, y) &\stackrel{\text{def}}{=} \texttt{penultV}(x) \land \texttt{LastH}(y) \\ \\ &\stackrel{\mathsf{H}}{\underset{\mathsf{V} \ \mathsf{V} \end{aligned}$$

Unbounded OCP

Unbounded OCP

$$H_O(x) \stackrel{\mathsf{def}}{=} H(x) \land \neg H(p(x))$$

Process		ISL?	AISL?
Bounded spread	(Bemba)	\checkmark	\checkmark
Bounded shift	(Rimi)	\checkmark	\checkmark
Unbounded shift	(Zigula)	X	 ✓
Unbounded OCP	(Arusa)	×	\checkmark
Unbounded spread	(Ndebele)	X	
Meussen's rule	(Shona)	X	

Unbounded spread





Unbounded spread

$$a_O(x, y) \stackrel{\text{def}}{=} a_I(x, y) \lor (\texttt{antepenultV}(x) \land H_I(y))$$



Unbounded spread

No QF statement can identify all intermediate vowels



Meussen's Rule

 Shona Odden (1986); Meyers (1987, 1997) /né-hóvé/ → [né-hòvè] /né-é-hóvé/ → [né-è-hóvé] /né-é-é-hóvé/ → [né-è-é-hòvè]

Meussen's Rule

 Need to pick out the set of even H's—this is well-known to be not definable even with (first-order) quantification (Thomas, 1982)

$$L_O \stackrel{\text{def}}{=} H_I(x) \wedge \text{even}(x)$$

Summary

Process		ISL?	AISL?
Bounded spread	(Bemba)	\checkmark	\checkmark
Bounded shift	(Rimi)	\checkmark	\checkmark
Unbounded shift	(Zigula)	X	\checkmark
Unbounded OCP	(Arusa)	×	 ✓
Unbounded spread	(Ndebele)	X	×
Meussen's rule	(Shona)	X	×

Discussion

Tone patterns include both ISL and non-ISL patterns Unbounded shift:

$$\vee \acute{\vee} \vee \vee \vee \vee \vee \mapsto \vee \vee \vee \vee \acute{\vee} \vee$$

▶ With AISL, we can capture some non-ISL patterns

Thus, ARs make some non-local patterns local

Discussion

Unbounded spread:

- One option for non-ISL/AISL processes is to further enrich the representations
- We might consider AR models with < instead of p (generalization of Strictly Piecewise (Heinz, 2010; Rogers et al., 2010))

Discussion

- Another option is to consider output-based locality
 - Johore Malay (Onn, 1980)
 /pəŋawasan/ → [pəŋāwãsan] 'supervision'
- Output SL functions have been characterized for strings in terms of formal language and automata theory (Chandlee et al., 2015)
- A logical characterization of OSL remains for future work.

Why do logical characterizations matter?

- Enable a rigorous, restrictive, and learnable (Chandlee and Heinz, 2018) definition of what it means to be "local" and "non-local".
- Directly extend these notions from strings to ARs.
- Logics are tightly connected to the complexity of functions (Filiot and Reynier, 2016).
- Computational complexity classes have been shown to capture the typology of spreading (Heinz and Lai, 2013).

Conclusion

- We *directly compared* different representations to better under how ARs can render non-local processes local.
- Given an input-based notion of locality, ARs capture some, but not all, patterns that are non-local over strings.
- In future, an output-based notion of locality may accommodate additional processes that are not AISL.

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